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(54) 5B6B coding for split channel transmission

(57) A data stream to be communicated over a plurality of channels is divided into blocks (A1,B1,C1,D1,A2,B2,...), and each successive block is transmitted along a different channel (A,B,C,D) on a cyclic basis. To reduce or eliminate the possibility of undetectable errors occurring owing to noise affecting all channels simultaneously and thereby corrupting data in several successive blocks propagating in parallel through the channels, the blocks in at least one channel (A,B) are offset in time relative to the blocks in another channel (C,D). In the case of four channels, such as four-conductor cable, the blocks on two channels are offset by half the length of a block relative to the blocks on the remaining two channels. A specific 5B/6B block code is also used in conjunction with a cyclic redundancy check algorithm to provide additional error detection capabilities.

Fig.8

5-bit data word	6-bit code
00000	001100 110011
00001	101100 101110
00010	001101 110101
00011	010101 001110
00100	001110 001011
00101	000111 100011
00110	100011 100110
00111	000110 111001
01000	101000 010110
01001	100100 011011
01010	101001 110110
01011	000101 100101
01100	001001 010110
01101	111000 011100
01110	011000 011001
01111	100001 011110
10000	110001 101010
10001	010100 110100
10010	110100 011100
10011	010010 101101
10100	
10101	
10110	
10111	
11000	
11001	
11010	
11011	
11100	
11101	
11110	
11111	

DescriptionTechnical Field

5 This invention relates to encoding and/or decoding data, for example in connection with communication of data over a plurality of channels, such as a cable having a plurality of conductors (for example a four-wire telephone cable).

Background Art

10 The spread of personal computers and workstations has led to the development of networks for interconnecting such equipment and common resources such as printers and data storage devices. More powerful and sophisticated computing equipment and programs have progressively become available, allowing the processing of data in larger and larger quantities, for example in the form of database information and graphic images. These developments have in turn placed increasing demands on the speed and capacity of networks.

15 Various new networking technologies have been proposed to cater for these demands. One such technology is the fibre distributed data interface (FDDI), which is based on the use of optical fibres and optical signals. However, practical experience has shown that although this technology can provide the required performance, it is relatively expensive, requiring the costly process of installing complete new networks of optical fibre, which is itself inherently expensive.

20 Accordingly attention has also been directed to the possibility of transferring data at high speed over existing wiring installations, thereby avoiding the cost of installing a new network and gaining additional return for the expense previously incurred in installing the existing wiring.

25 One possible technique along these lines involves the use of unshielded twisted-pair (UTP) telephone cables of the kind already used for lower-speed local-area networks. In this technique, described for example in US Patent 5 119 402, the required high data rate is achieved by transferring the data over multiple conductors, so that different portions of the data are transmitted simultaneously over respective conductors.

It is an object of this invention to provide methods for encoding and decoding, and associated coders and decoders, for use in these and other systems.

Disclosure of Invention

30 According to one aspect of this invention there is provided a method of encoding 5-bit data words as 6-bit code words for use with a cyclic redundancy check algorithm based upon the polynomial

$$g(x)=x^{32}+x^{26}+x^{23}+x^{22}+x^{16}+x^{12}+x^{11}+x^{10}+x^8+x^7+x^5+x^4+x^2+x+1$$

35 wherein 6-bit code words are selected in accordance with the values of the 5-bit data words and a preselected one of the tables shown in Figures 3 and 8, or an invariance transformation thereof, code word selection in the case of data words having two possible code words being of the code word from a column of the preselected table alternate to the column from which a code word was selected upon a preceding occurrence of any of such data words. In a complementary method of decoding 5-bit data words from 6-bit code words, 5-bit data words are selected in accordance with the values of the 6-bit code words and a preselected one of the tables shown in Figures 3 and 8, or an invariance transformation thereof. Corresponding coders and decoders are also provided.

Brief Description of Drawings

45 Methods and apparatus in accordance with this invention for encoding and decoding data for communication using four-conductor unshielded twisted-pair cable will now be described, by way of example, with reference to the accompanying drawings, in which:

- 50 Figure 1 shows the manner in which data is formatted for communication over the cable;
- Figure 2 illustrates the effect of noise on data communicated over the cable;
- Figure 3 is a table of five-bit data values and corresponding six-bit code values;
- Figure 4 illustrates reduction possible in the effect of noise on data communicated over the cable;
- Figure 5 is a block schematic diagram of part of a transmitter incorporating the present invention;
- 55 Figure 6 is a block schematic diagram of part of a receiver incorporating the present invention;
- Figure 7 is a flow diagram of a method for encoding data blocks according to the table in Figure 3; and
- Figure 8 is an alternative table of five-bit data values and corresponding six-bit code values.

Best Mode for Carrying Out the Invention, & Industrial Applicability

The present invention may be used, for example, in circumstances where a stream of data is communicated over a plurality of channels, successive portions of the data stream being communicated simultaneously over different
 5 respective channels in order to obtain a higher bandwidth than would be possible if all the data were transmitted over a single such channel. For convenience the invention will be described in the context of transmission of binary data over a cable having four channels or conductors (e.g. four pairs of twisted wires). However, the invention is not limited to this particular number of channels nor to this type of channel. In practice the cable would, for example, form part of a network connecting many stations or nodes, such as personal computers, workstations, multi-user computers, printers or
 10 data storage units. Circuit devices associated with these stations would provide the necessary functions for assembling data and network operating information into frames or packets for transmission, for controlling access to the network and for transmitting and receiving physical signals on the cable (for example by differential signalling in the case of twisted-pair conductors). The present invention is independent of the particular details of these functions and may for example be implemented in conjunction with existing network technologies; since such technologies already incorporate
 15 known techniques for providing these functions, and the functions form no part of the present invention, they will not be described here.

Referring to Figures 1 and 2, a data frame intended to be communicated over a four-conductor cable is shown schematically at 10. This frame comprises: a binary digital message 12 to be transferred, starting with the leftmost bit as shown in Figure 1, between stations on the network; and an associated thirty-two bit CRC block 14 containing check
 20 data derived from the message 12 in known manner in accordance with a predetermined cyclic redundancy check (CRC) algorithm. In the present example it is assumed that the CRC value is derived from the message using a polynomial of degree thirty-two, such as

$$g(x) = x^{32} + x^{26} + x^{23} + x^{22} + x^{16} + x^{12} + x^{11} + x^{10} + x^8 + x^7 + x^5 + x^4 + x^2 + x + 1.$$

For transmission over the four-conductor cable the data frame 10 is split into consecutive blocks of five bits each, and the blocks are distributed among the four conductors (herein labelled A to D) on a cyclic basis and starting with the block containing the leftmost bit of the data frame. Thus this first block, labelled A1 in Figure 1, is transmitted via conductor A, the next block (B1) via conductor B, the third block (C1) via conductor C and the fourth (D1) via conductor D.
 30 The cycle then repeats, with conductor A being used again, for the fifth block (A2), and so on.

Prior to transmission the five-bit data blocks are encoded by a 5-to-6 bit encoder 16 (Figure 2) into six-bit values according to a substitution table, to provide a measure of inherent error detection. The encoding substitutions may be as shown in Figure 3. These particular substitutions are selected in part to maintain d.c. balance on each conductor, by ensuring that after transmission of each coded data block the accumulated totals of binary ones and zeroes differ by no
 35 more than two. Thus, in the example shown in Figure 3, twenty substitutions are assigned unique codes comprising respective ones of the twenty six-bit values which contain three binary zeroes and three binary ones. The remaining twelve data blocks are each assigned two possible six-bit code values, one containing two binary zeroes and the other containing four. The encoding is implemented so that on the first occasion one of these twelve data blocks occurs for transmission along a particular conductor, the corresponding two-zero encoding (for example) is selected; on the next
 40 occasion that any of these twelve data blocks occurs for transmission along that same conductor, the corresponding four-zero encoding is used; thereafter the use of the two-zero and four-zero encodings continues to alternate for each occurrence of any of these twelve data blocks in respect of that conductor. Thus for each conductor the numbers of two-zero and four-zero six-bit codes will differ at most by one, maintaining an average of three binary zeroes per six-bit code and providing the desired d.c. balance. In Figure 1 illustrative five-bit data blocks are indicated in bold characters, and
 45 corresponding six-bit code values are indicated below them in normal weight characters.

The particular set of substitutions given in Figure 3 is illustrative only; different arrangements of five-bit data values and six-bit code values may be assigned as desired, as described below with reference to Figure 8.

After encoding, the six-bit code values are distributed or 'de-multiplexed' by a de-multiplexer 18 among the four conductors A to D on a cyclic basis as described above. Thus, as shown in Figure 2, conductor A will carry encoded data
 50 blocks A1, A2, A3, etc. in succession (leftmost bit of each first), conductor B will carry encoded data blocks B1, B2, B3, etc., and likewise for conductors C and D.

Figure 2 also illustrates the potential effects of bursts of electrical noise on the data carried by the conductors A to D. Typically such a noise burst can interfere with the data for a period as long as the interval occupied by four of the six bits of an encoded data block, as indicated by the dashed lines. Thus, if a noise burst happens to commence at or just
 55 after the start of a data block, as illustrated by the noise burst 1 in Figure 2, the encoded data block on each of the four conductors may be corrupted. As a result of the encoding process, such corruption can completely alter the value obtained upon decoding, even though only some of the bits of the encoded data block are affected. The corrupted encoded data blocks correspond to four successive five-bit data values of the original data frame. Thus up to twenty

successive bits of the data frame may be corrupted. Since the CRC algorithm is based on a polynomial of degree thirty-two, errors involving twenty successive bits are always detected.

However, it is also possible for noise to affect two successive data blocks on each conductor, as illustrated by the noise burst 2 in Figure 2. This noise burst straddles the end of one data block and the start of the next data block on each conductor. Thus, the final two bits of encoded data block A1 (01) and the first two bits of data block A2 (11) may be corrupted. Consequently a total of eight successive data blocks (forty successive bits) of the original data frame may be corrupted. A CRC algorithm based on a polynomial of degree thirty-two cannot be guaranteed to detect errors affecting this many data bits, so there is a possibility of undetected errors occurring. Although a longer CRC polynomial could be used, this would impose an additional and undesirable processing burden on the equipment included in the network.

Figure 4 illustrates a solution to this problem. In contrast to Figure 2, the encoded data blocks in Figure 4 are transmitted on different ones of the four conductors temporally offset from one another. In the specific example shown in Figure 4, data blocks are transmitted on conductors A and B coincidentally (or, in terms of overall blocks, in phase with each other), and offset from (or out of phase with) the data blocks transmitted coincidentally with one another on conductors C and D. The offset in this case is equal to half the length of the encoded data blocks. Thus encoded data blocks on conductors C and D commence half-way through the transmission of encoded data blocks on conductors A and B.

With this arrangement it is still possible for a burst of noise arbitrarily corrupting four successive code bits on all four conductors simultaneously to straddle two encoded data blocks on each of two of the conductors (see the noise burst 3 in Figure 4). However, by virtue of the offset timing of the data blocks on the other two conductors, this noise burst can only affect one data block at most on each of the other two conductors. If the noise burst 3 were to occur two code bits later, as it would have to do to affect encoded blocks C3 and D3 at all, it would no longer affect blocks A2 and B2; although two blocks would now be affected on conductors C and D, only one block would be affected on each of A and B. As a result, the maximum number of consecutive data blocks which can be corrupted is reduced to six, corresponding to thirty successive bits of the original data frame. Corruption even of this many data bits can be reliably detected by a CRC algorithm using a degree thirty-two polynomial.

The required offset in the transmission times of data blocks on the different conductors may be provided in various ways. One possible approach involves the use of shift registers as shown in Figures 5 and 6.

Referring to Figure 5, data for transmission is supplied to a circuit 20 which performs the steps of calculating the CRC value and appending it to the message data, and encoding successive blocks of five data bits to derive six-bit encoded blocks as described above. The encoded data blocks are supplied to a media access control (MAC) circuit 22 which coordinates access to the network cable in accordance with a predetermined protocol to ensure efficient use of the communication bandwidth provided by the cable. The MAC circuit 22 passes the encoded blocks to a four-way 'de-multiplexer' 24, which distributes the blocks cyclically among four outputs *a* to *d*, and generates a clock signal on a line 26 which is synchronized with the digital signals appearing on these outputs. The circuit 24 is described herein as a de-multiplexer even though its input signal is not a multiplex signal in the conventional sense, since its function is essentially that of de-multiplexing: cyclically routing successive portions of the incoming signal to respective ones of its outputs according to a predetermined pattern.

The outputs *a* and *b* of the de-multiplexer 24 are coupled directly to the conductors A and B; however, the outputs *c* and *d* are coupled to the inputs of respective three-bit shift registers 28 and 30. These shift registers each receive the clock signal on the line 26, and their outputs are coupled to the conductors C and D. Thus the signals on the outputs *c* and *d* of the de-multiplexer 24 are actually propagated over the conductors C and D with a time delay of three bit periods (half the duration of a complete encoded data block) relative to the signals on the conductors A and B.

Referring to Figure 6, a pulse signal on the conductor A is coupled to the input of a phase-locked loop (PLL) 32, which generates a clock signal on a line 34 in synchronism with the incoming pulse signal. This signal is also supplied to the input of a threshold circuit 36 which restores voltage levels occurring in the pulse signal to the correct values indicative of binary zero and one, by comparison with predetermined threshold magnitudes. The restored pulse signal is coupled to the input of a sample and decision circuit 38 which samples the signal in synchronism with the clock signal on the line 34 and generates an output digital signal phase-locked to the clock signal and having a binary value dependent on the value of the sample.

The output of the sample and decision circuit 38 is connected to the input of a three-bit shift register 40 which also receives the clock signal on the line 34 and has its output coupled to an input *a* of a de-skew and multiplexing circuit 42 which also receives the associated clock signal. The pulse signal on the conductor B is treated in a similar manner to produce an output digital signal which is coupled via a three-bit shift register 44 to an input *b* of the circuit 42. This signal is phase-locked to a respective clock signal which is also supplied to the circuit 42.

Similar processing is applied to the pulse signals on the conductors C and D by respective circuitry associated with those conductors to produce associated output digital signals and clock signals. However, these two output digital signals are coupled directly from the associated sample and decision circuits to inputs *c* and *d* of the de-skew and multiplexing circuit 42, without traversing any intervening shift register.

The de-skew and multiplexing circuit 42 applies minor timing corrections to the four output digital signals, to com-

pensate for any minor timing skew which may have arisen between them during propagation along the four-conductor cable, and then 'multiplexes' the four signals by coupling blocks of six bits from each input in turn to its output to reconstruct the original encoded data stream.

This encoded data stream is fed to a circuit 46 which decodes each six-bit encoded block to derive the corresponding five-bit data block. If any six-bit encoded value is encountered which is invalid or which has an incorrect number of binary zero bits the entire data frame is rejected. Otherwise the circuit 46 assembles the complete data frame and recalculates the CRC value for comparison with the transmitted CRC value.

The inclusion of the shift registers 40 and 44 at the reception end of the signal path for the signals on the conductors A and B introduces a three-bit delay which matches that provided by the shift registers 28 and 30 in Figure 5, so that the signals on all four conductors experience the same total delay from this source. However, during actual propagation along the conductors the data blocks on conductors C and D are offset relative to those on conductors A and B, with the consequent advantages described above.

It should be noted that the communication of five-bit data blocks and six-bit encoded blocks over four channels comprising wire conductors in a cable as described above is purely illustrative; other lengths of data block and numbers and types of channels may be used. Likewise the three-bit offset between data blocks on different conductors is an optimum value related to the particular block length chosen for the purposes of description. The actual offset in any particular case may differ from this amount, and need not be half the block length.

The arrangement described above ensures that all errors caused by four code bit noise pulses will affect no more than thirty consecutive data bits (in the particular example given), and thus will be detectable with a thirty-two bit CRC algorithm. However, it is possible, by appropriate choice of the correspondence between (five-bit) data values and (six-bit) code values, to avoid the occurrence of errors which affect larger numbers of consecutive data bits and which would not generally be detectable without resorting to a CRC algorithm based on a higher degree polynomial.

Thus, for example, to avoid undetectable errors caused by noise bursts more than four code bits in duration and affecting up to forty consecutive data bits, with a CRC polynomial of degree thirty-two, a list may be prepared of all possible forty-bit errors which are not detectable by the CRC algorithm. In this context 'error' means the result of an exclusive-OR operation between transmitted data and corrupted received data. For any particular CRC polynomial of degree thirty-two there are two hundred and fifty-five such error values.

An analysis may then be performed of possible forty-bit errors which can arise from noise affecting data being communicated, for different choices of five-bit to six-bit encodings (it will be apparent to those skilled in the art that there is an extremely large population of such encodings). Typically it will be found that only some of the listed undetectable errors will arise for any particular choice of encoding. However, for at least some CRC polynomials there exist certain encodings which, however corrupted by a noise burst up to six code bits in duration (and involving up to a maximum of forty consecutive data bits), do not give rise to any of these two hundred and fifty-five undetectable errors for any data value. Identification and use of such an encoding in conjunction with the offsetting of data blocks on different channels as described above will ensure that all errors arising from noise bursts up to six code bits in duration will be reliably detected. The encoding given in Figure 3 is an example of an encoding having this property.

Figure 8 is an example of another encoding having similar properties, but which, used in conjunction with the offsetting of data blocks on different channels, will ensure that all errors arising from noise bursts up to seven code bits in duration (and thus arbitrarily corrupting seven successive code bits on all four conductors simultaneously) will be reliably detected. Errors arising from a combination of noise and signal skew as between the different conductors A to D, and together having a duration of up to seven code bits, will likewise be reliably detected.

Figure 7 shows a flow diagram of a method for implementing such an encoding. Referring to Figure 7, at step 100 a counter K is incremented according to the relationship

$$K = (K + 1) \text{ modulo } 4$$

so that the counter cyclically takes on the values zero to three inclusive. The purpose of this counter is to keep track of which conductor the current (encoded) data value will be transmitted along. At step 102 the data value to be encoded is tested to check whether it has one or two corresponding six-bit code values. If there is a unique corresponding code value, the procedure obtains that value from a look-up table at step 104 and exits. Otherwise the procedure advances to step 106, where one of four boolean flags, selected in accordance with the current value of the counter K, is tested. If the flag is 'true', the procedure selects the code value containing only two binary zeroes, at step 108; if the flag is 'false' the code value containing four zeroes is selected, at step 110. In either case the procedure then inverts the value of the flag at step 112 before exiting.

A broadly similar procedure to that of Figure 7 may be used by the circuit 46 for decoding. However, instead of obtaining a code value at step 104, the circuit 46 would check that the received code value is valid, and then obtain the corresponding data value. Likewise, at steps 108 and 110, the circuit 46 would check that the received code value is valid and has the expected number of binary zero bits, and then obtain the required data value. If any of these checks

failed, the circuit 46 would determine that an error had occurred during transmission of the data frame.

It will be obvious to a person skilled in the art, given the encodings of Figure 3 or Figure 8, that additional encodings having the same properties may be derived therefrom by simple invariance transformations. As an example of such a transformation, the encoding of Figure 8 may be transformed to an equivalent form by inverting the bit values of all the six-bit code values in the right-hand part (three columns) of the Figure; this process is equivalent to an exclusive-OR operation with a constant binary value of 111111. Further equivalent encodings result from an exclusive-OR operation performed with any constant five-bit binary value on all the five-bit data values in the left-hand part of the Figure; such an encoding may be further transformed by inverting the bit values of the six-bit code values, as described above, to produce another equivalent encoding.

Claims

1. A method of encoding 5-bit data words as 6-bit code words for use with a cyclic redundancy check algorithm based upon the polynomial

$$g(x)=x^{32}+x^{26}+x^{23}+x^{22}+x^{16}+x^{12}+x^{11}+x^{10}+x^8+x^7+x^5+x^4+x^2+x+1$$

wherein 6-bit code words are selected in accordance with the values of the 5-bit data words and a preselected one of first and second tables set out below, or an invariance transformation thereof, code word selection in the case of data words having two possible code words being of the code word from a column of the preselected table alternate to the column from which a code word was selected upon a preceding occurrence of any of such data words:

First Table

	<u>5-bit data word</u>	<u>6-bit code word</u>	
5	00000	111001	000110
	00001		001110
	00010		110010
10	00011		000111
	00100		100110
	00101		010011
15	00110	011110	100001
	00111	100111	011000
20	01000		110100
	01001		010110
	01010	111010	000101
25	01011		100011
	01100		110001
	01101	110110	001001
	01110		011010
30	01111		010101
35	10000	101011	010100
	10001	011011	100100
	10010		100101
	10011		101010
40	10100		001011
	10101		101001

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	10110	010111	101000
	10111	110101	001010
5			
	11000	011001	
	11001	101100	
10	11010	101101	010010
	11011	011100	
	11100	101110	100010
	11101	110011	001100
15	11110	001101	
	11111	111000	

Second Table

20	<u>5-bit data word</u>	<u>6-bit code word</u>	
	00000	001100	110011
	00001	101100	
25	00010	100010	101110
	00011	001101	
	00100	001010	110101
30	00101	010101	
	00110	001110	
	00111	001011	
35	01000	000111	
	01001	100011	
	01010	100110	
40	01011	000110	111001
	01100	101000	010111
	01101	011010	
	01110	100100	011011
45	01111	101001	
	10000	000101	111010
50	10001	100101	
	10010	001001	110110
	10011	010110	

	10100		111000
	10101	011000	100111
5	10110		011001
	10111	100001	011110
10	11000		110001
	11001		101010
	11010	010100	101011
15	11011		110100
	11100		011100
	11101		010011
20	11110	010010	101101
	11111		110010

2. A coder for encoding 5-bit data words as 6-bit code words for use with a cyclic redundancy check algorithm based upon the polynomial

$$g(x) = x^{32} + x^{26} + x^{23} + x^{22} + x^{16} + x^{12} + x^{11} + x^{10} + x^8 + x^7 + x^5 + x^4 + x^2 + x + 1$$

comprising means for receiving 5-bit data words and for selecting 6-bit code words in accordance with the values of the 5-bit data words and a preselected one of first and second tables as set out below, or an invariance transformation thereof, and including means for selecting, in the case of data words having two possible code words, the code word from a column of the preselected table alternate to the column from which a code word was selected upon a preceding occurrence of any of such data words:

First Table

	<u>5-bit data word</u>	<u>6-bit code word</u>	
40	00000	111001	000110
	00001		001110
	00010		110010
45	00011		000111
	00100		100110
	00101		010011
50	00110	011110	100001
	00111	100111	011000
55	01000		110100
	01001		010110

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	01010	111010	000101
5	01011		100011
	01100		110001
	01101	110110	001001
10	01110		011010
	01111		010101
	10000	101011	010100
15	10001	011011	100100
	10010		100101
	10011		101010
20	10100		001011
	10101		101001
	10110	010111	101000
25	10111	110101	001010
	11000		011001
	11001		101100
30	11010	101101	010010
	11011		011100
	11100	101110	100010
35	11101	110011	001100
	11110		001101
	11111		111000

Second Table

	<u>5-bit data word</u>	<u>6-bit code word</u>	
	00000	001100	110011
45	00001		101100
	00010	100010	101110
	00011		001101
50	00100	001010	110101
	00101		010101
	00110		001110
55	00111		001011

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	01000		000111
	01001		100011
5	01010		100110
	01011	000110	111001
	01100	101000	010111
10	01101		011010
	01110	100100	011011
	01111		101001
15			
	10000	000101	111010
	10001		100101
20	10010	001001	110110
	10011		010110
	10100		111000
25	10101	011000	100111
	10110		011001
	10111	100001	011110
30			
	11000		110001
	11001		101010
35	11010	010100	101011
	11011		110100
	11100		011100
	11101		010011
40	11110	010010	101101
	11111		110010

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3. A method of decoding 5-bit data words from 6-bit code words, wherein 5-bit data words are selected in accordance with the values of the 6-bit code words and a preselected one of first and second tables set out below, or an invar-
- 50 iance transformation thereof:

55

First Table

5-bit data word

6-bit code word

5	00000	111001	000110
	00001		001110
	00010		110010
10	00011		000111

15

20

25

30

35

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45

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	00100		100110
	00101		010011
5	00110	011110	100001
	00111	100111	011000
10	01000		110100
	01001		010110
	01010	111010	000101
	01011		100011
15	01100		110001
	01101	110110	001001
	01110		011010
20	01111		010101
	10000	101011	010100
	10001	011011	100100
25	10010		100101
	10011		101010
	10100		001011
	10101		101001
30	10110	010111	101000
	10111	110101	001010
35	11000		011001
	11001		101100
	11010	101101	010010
	11011		011100
40	11100	101110	100010
	11101	110011	001100
	11110		001101
45	11111		111000

Second Table

	<u>5-bit data word</u>	<u>6-bit code word</u>
50	00000	001100 110011
	00001	101100

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	00010	100010	101110
	00011	001101	
5	00100	001010	110101
	00101	010101	
	00110	001110	
10	00111	001011	
	01000	000111	
15	01001	100011	
	01010	100110	
	01011	000110	111001
20	01100	101000	010111
	01101	011010	
	01110	100100	011011
25	01111	101001	
	10000	000101	111010
	10001	100101	
30	10010	001001	110110
	10011	010110	
	10100	111000	
35	10101	011000	100111
	10110	011001	
	10111	100001	011110
40			
	11000	110001	
	11001	101010	
45	11010	010100	101011
	11011	110100	
	11100	011100	
	11101	010011	
50	11110	010010	101101
	11111	110010	
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4. A decoder for decoding 5-bit data words from 6-bit code words, comprising means for receiving 6-bit code words and for selecting 5-bit data words in accordance with the values of the 6-bit code words and a preselected one of

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first and second tables set out below, or an invariance transformation thereof:

First Table

	<u>5-bit data word</u>	<u>6-bit code word</u>	
5	00000	111001	000110
	00001		001110
	00010		110010
10	00011		000111
	00100		100110
	00101		010011
15	00110	011110	100001
	00111	100111	011000
20	01000		110100
	01001		010110
	01010	111010	000101
25	01011		100011
	01100		110001
	01101	110110	001001
	01110		011010
30	01111		010101
35	10000	101011	010100
	10001	011011	100100
	10010		100101
	10011		101010
40	10100		001011
	10101		101001
	10110	010111	101000
45	10111	110101	001010
	11000		011001
	11001		101100
50	11010	101101	010010
	11011		011100
55	11100	101110	100010

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	11101	110011	001100
5	11110		001101
	11111		111000

Second Table

10	<u>5-bit data word</u>	<u>6-bit code word</u>	
	00000	001100	110011
	00001		101100
15	00010	100010	101110
	00011		001101
	00100	001010	110101
	00101		010101
20	00110		001110
	00111		001011
25	01000		000111
	01001		100011
	01010		100110
	01011	000110	111001
30	01100	101000	010111
	01101		011010
	01110	100100	011011
35	01111		101001
40	10000	000101	111010
	10001		100101
	10010	001001	110110
	10011		010110
	10100		111000
45	10101	011000	100111
	10110		011001
	10111	100001	011110
50	11000		110001
	11001		101010
55	11010	010100	101011

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11011	110100
11100	011100
11101	010011
11110	010010 101101
11111	110010

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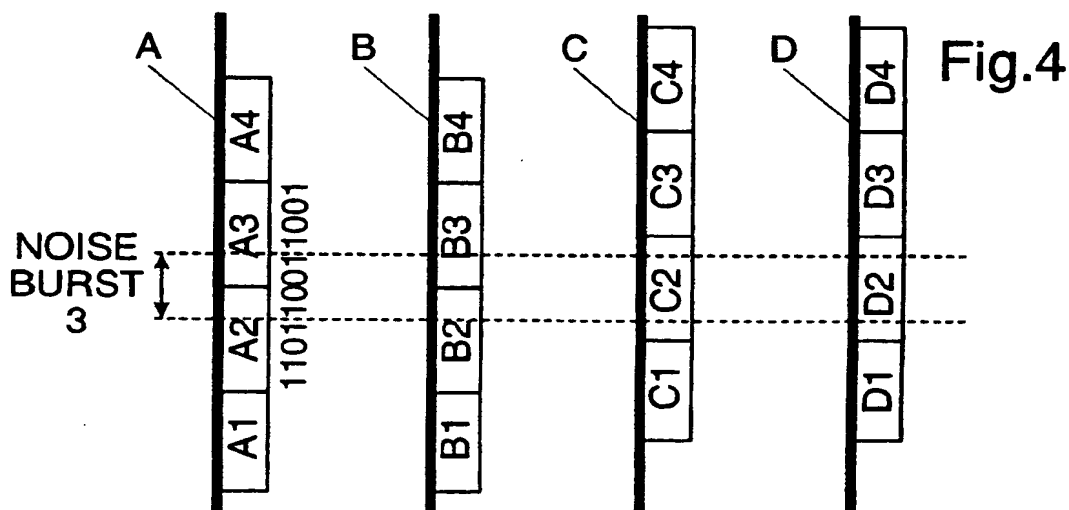
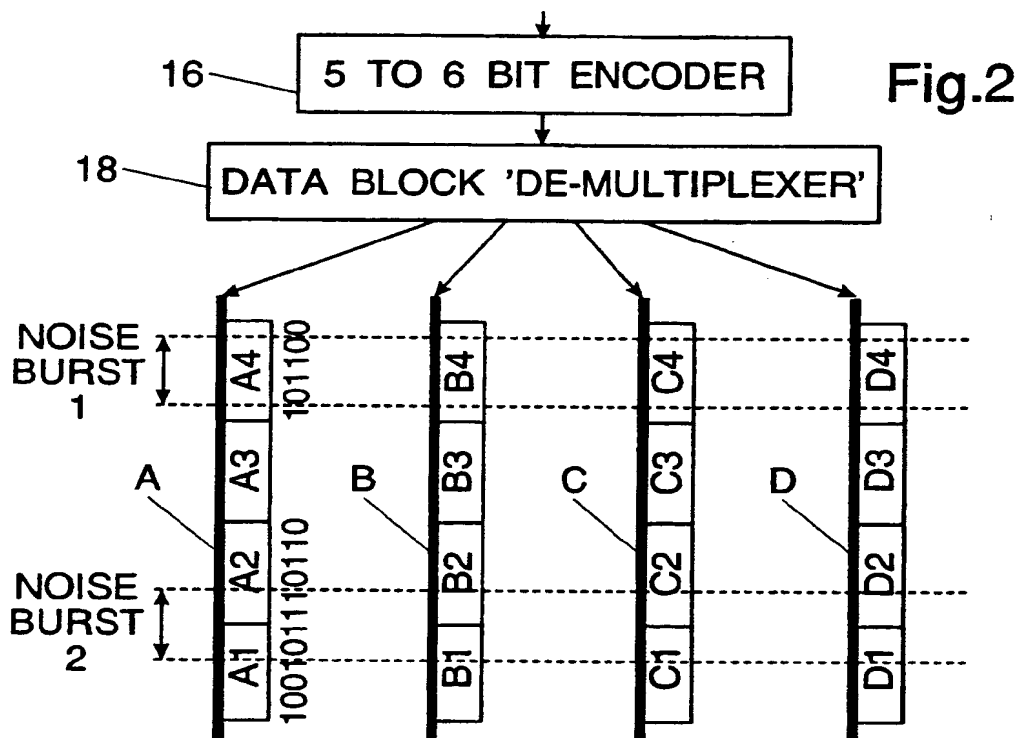
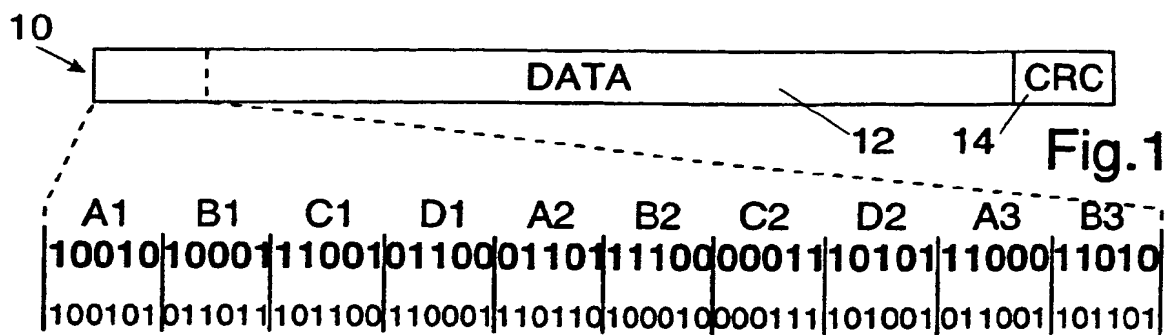


Fig.3

<u>5-bit data word</u>		<u>6-bit code</u>	
00000	111001		000110
00001		001110	
00010		110010	
00011		000111	
00100		100110	
00101		010011	
00110	011110		100001
00111	100111		011000
01000		110100	
01001		010110	
01010	111010		000101
01011		100011	
01100		110001	
01101	110110		001001
01110		011010	
01111		010101	
10000	101011		010100
10001	011011		100100
10010		100101	
10011		101010	
10100		001011	
10101		101001	
10110	010111		101000
10111	110101		001010
11000		011001	
11001		101100	
11010	101101		010010
11011		011100	
11100	101110		100010
11101	110011		001100
11110		001101	
11111		111000	

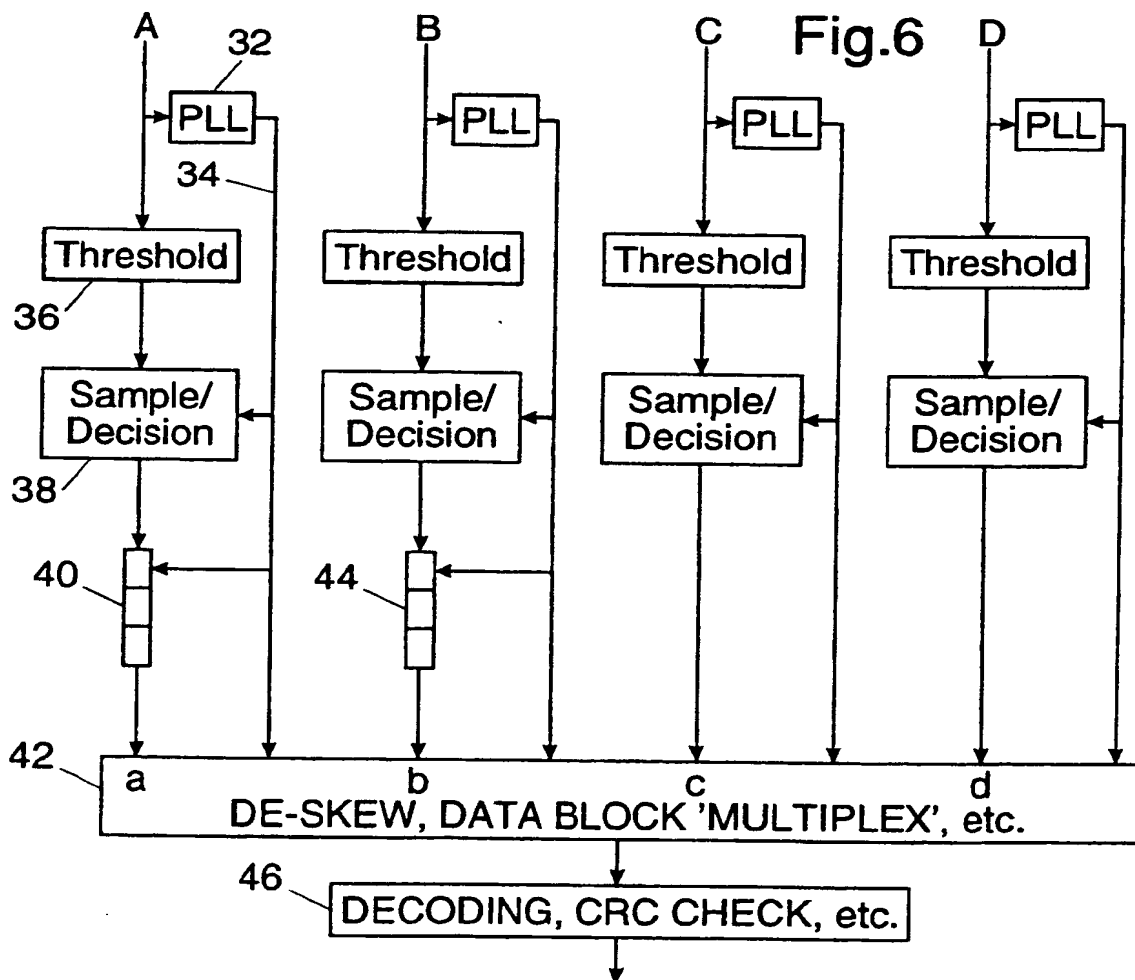
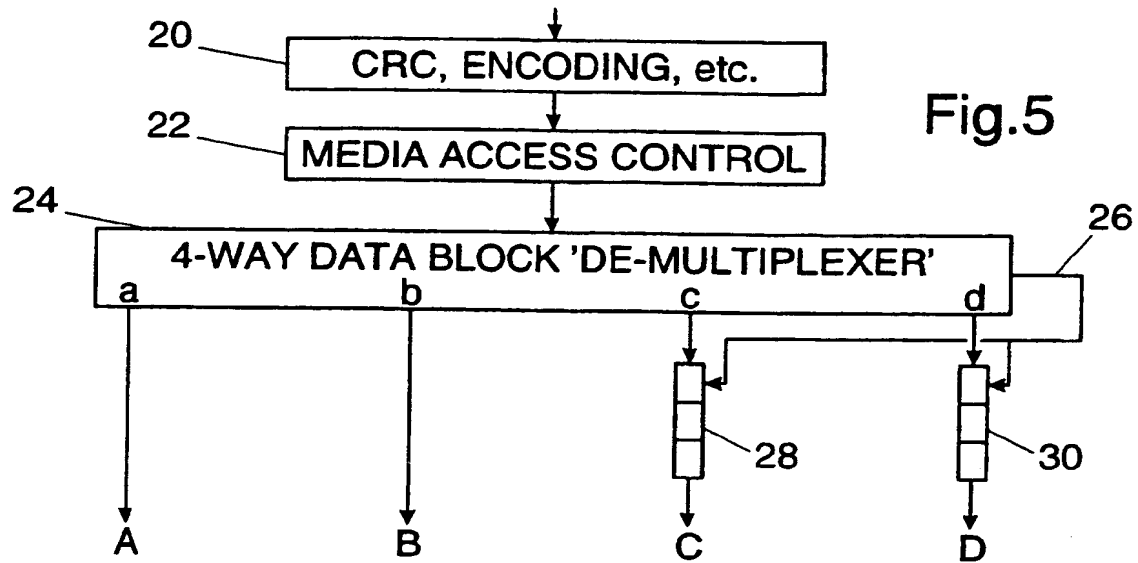


FIG.7

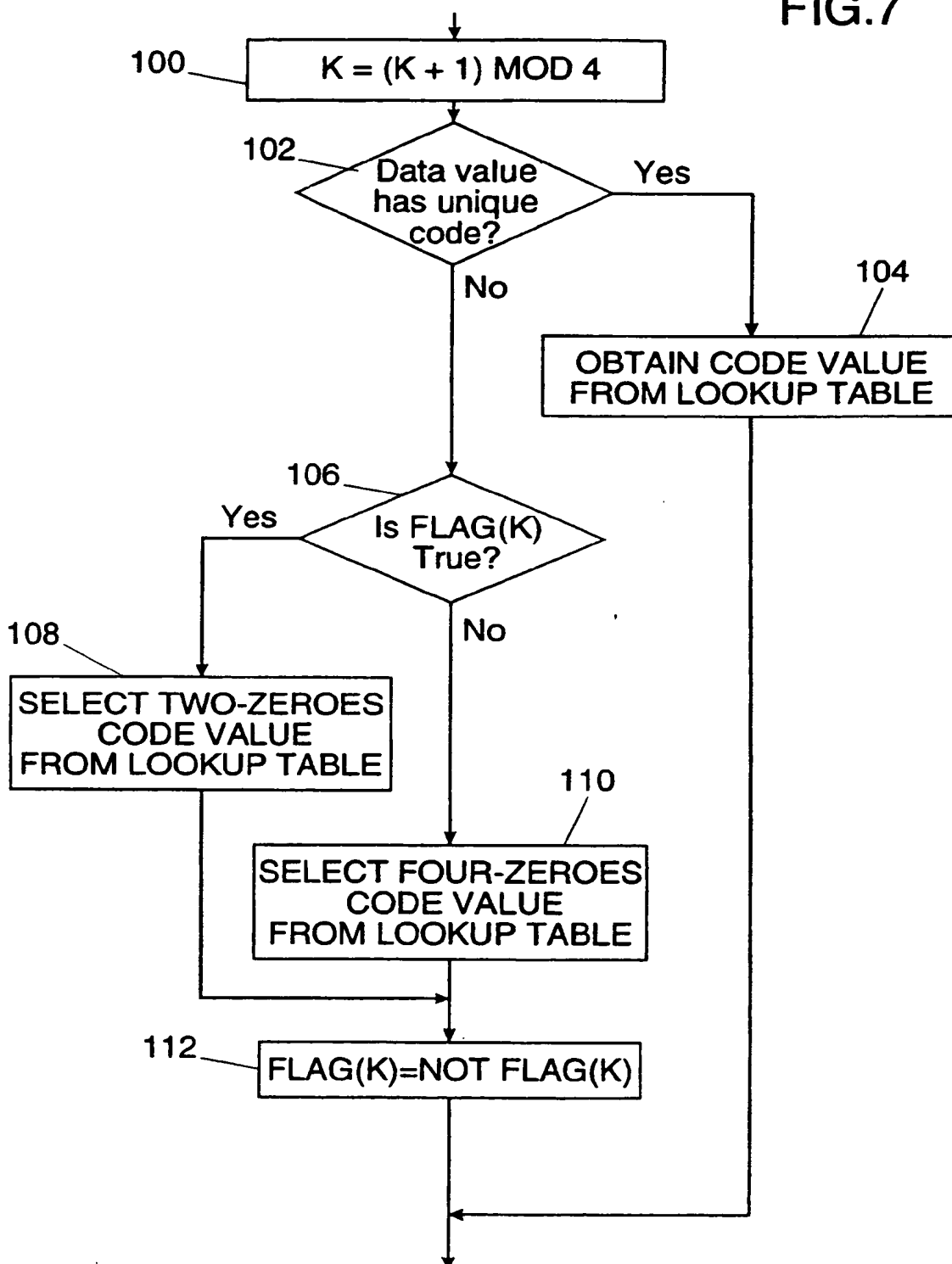


Fig.8

<u>5-bit data word</u>	<u>6-bit code</u>	
00000	001100	110011
00001		101100
00010	100010	101110
00011		001101
00100	001010	110101
00101		010101
00110		001110
00111		001011
01000		000111
01001		100011
01010		100110
01011	000110	111001
01100	101000	010111
01101		011010
01110	100100	011011
01111		101001
10000	000101	111010
10001		100101
10010	001001	110110
10011		010110
10100		111000
10101	011000	100111
10110		011001
10111	100001	011110
11000		110001
11001		101010
11010	010100	101011
11011		110100
11100		011100
11101		010011
11110	010010	101101
11111		110010

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(54) 5B6B coding for split channel transmission

(57) A data stream to be communicated over a plurality of channels is divided into blocks (A1,B1,C1,D1,A2,B2,...), and each successive block is transmitted along a different channel (A,B,C,D) on a cyclic basis. To reduce or eliminate the possibility of undetectable errors occurring owing to noise affecting all channels simultaneously and thereby corrupting data in several successive blocks propagating in parallel through the channels, the blocks in at least one channel (A,B) are offset in time relative to the blocks in another channel (C,D). In the case of four channels, such as four-conductor cable, the blocks on two channels are offset by half the length of a block relative to the blocks on the remaining two channels. A specific 5B/6B block code is also used in conjunction with a cyclic redundancy check algorithm to provide additional error detection capabilities.

Fig.8

5-bit data word	6-bit code	6-bit code
00000	001100	110011
00001	101100	101110
00010	001101	110101
00011	010101	001110
00100	001110	001011
00101	000111	100011
00110	100011	100110
00111	000110	111001
01000	101000	010111
01001	001101	111010
01010	100100	101011
01011	101001	110110
01100	001001	100101
01101	010100	101011
01110	011000	110011
01111	010001	011110
10000	110001	101010
10001	101010	101011
10010	110100	011100
10011	011100	010011
10100	110011	101101
10101	011101	110101
10110	110101	101101
10111	110110	110101



European Patent
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EUROPEAN SEARCH REPORT

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DOCUMENTS CONSIDERED TO BE RELEVANT			
Category	Citation of document with indication, where appropriate, of relevant passages	Relevant to claim	CLASSIFICATION OF THE APPLICATION (Int.Cl.6)
X	EP 0 484 946 A (FUJITSU) 13 May 1992 * figure 1 *	1-4	H04L25/49 H04L25/14 H04L1/00
X	US 3 631 471 A (GRIFFITHS) 28 December 1971 * table 1 *	1-4	
X	--- PETROVIC: "586B optical fibre line code bearing auxilliary signals" ELECTRONICS LETTERS, vol. 24, no. 5, 3 March 1988, pages 274-275, XP002082917 ENAGE, GB * table 1 *	1-4	
A	--- US 4 712 215 A (JOSHI, IYER) 8 December 1987 * column 1, line 53 - line 64 *	1,2	
X,P	--- EP 0 556 981 A (ADVANCED MICRO DEVICES) 25 August 1993 * table 1 *	1-4	
	-----		TECHNICAL FIELDS SEARCHED (Int.Cl.6)
			H04L
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